

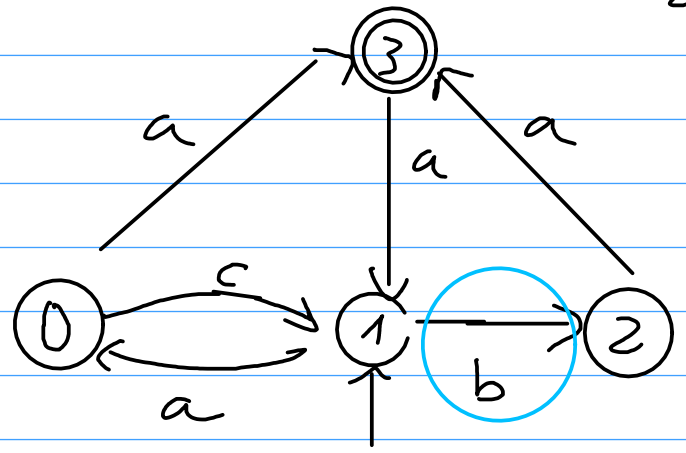
Glushkov $L = (a + \varepsilon) (b + aa)^* (b + \varepsilon)$

Linearisation $x_2 \rightarrow x_1$ $\varepsilon \rightarrow \text{eps or } \varepsilon$

$(x_1 + \varepsilon) (x_2 + x_3 x_4)^* (x_5 + \varepsilon)$

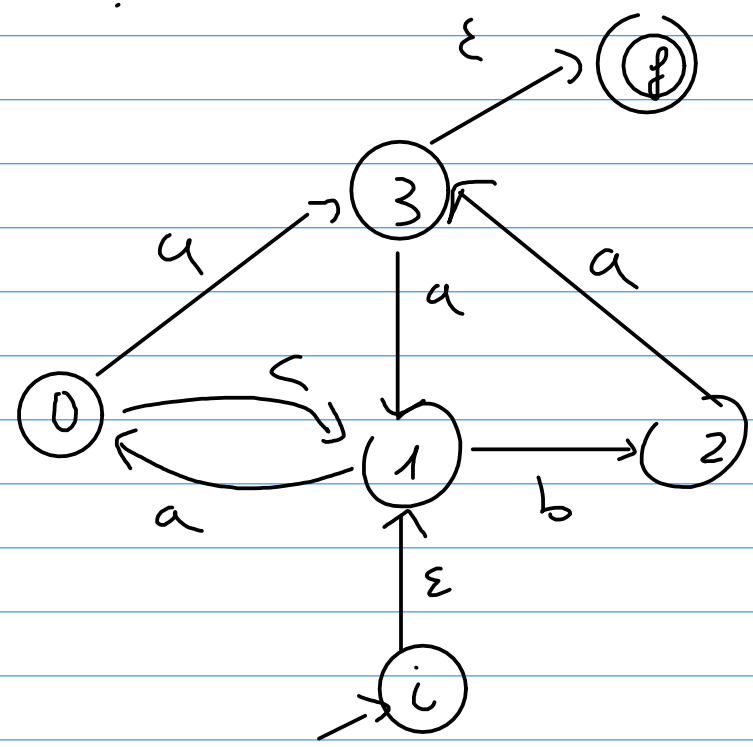
	Success.		a	b	
0	1, 2, 3, 5, fin.	0	1, 3	2, 5	↔
1	2, 3, 5, fin	1	3	2, 5	↑
2	2, 3, 5, fin	2	3	2, 5	↑
3	4	3	4		
4	2, 3, 5, fin	4	3	2, 5	↑
5	fin	5			↑

Brzozowski : t_2 du TD4 Exo 1

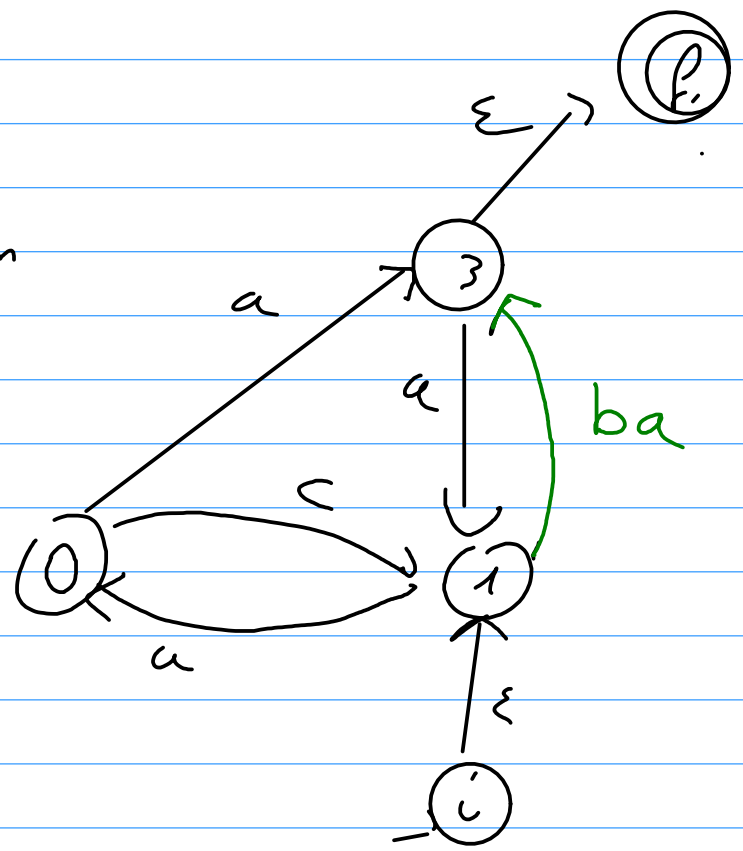


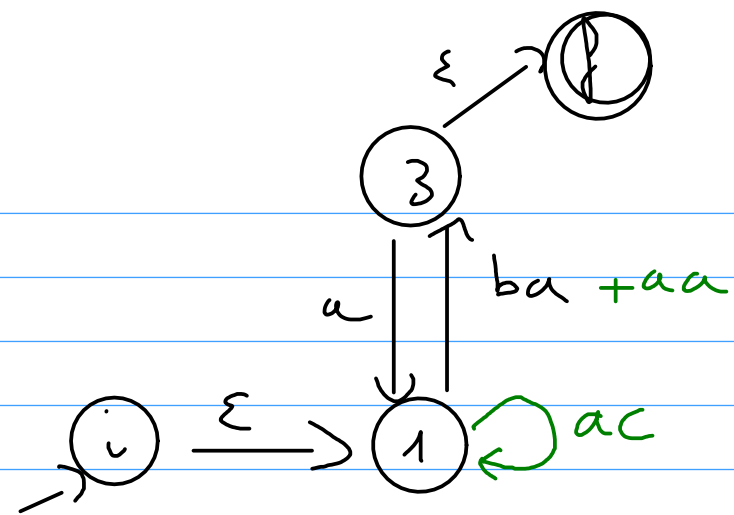
ordre d'élimination :

(2) , (0) , 3 , 1

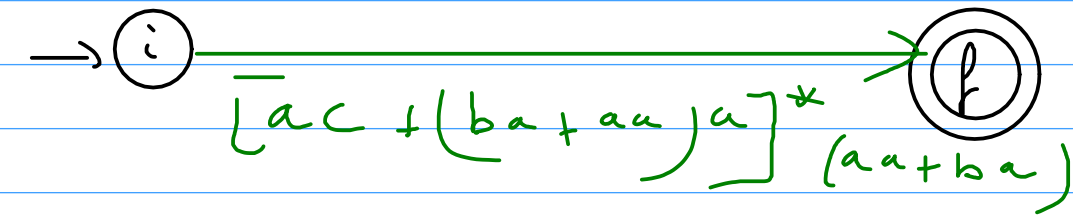
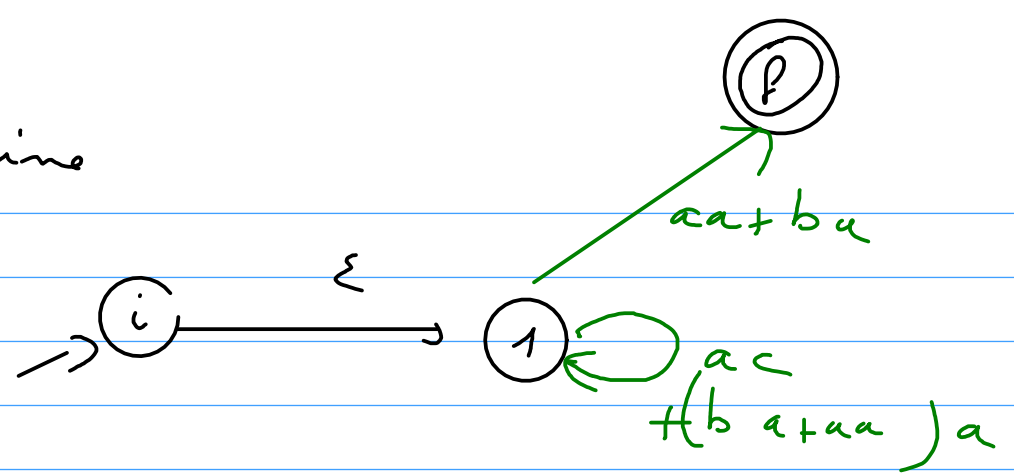


élimination
de 2





Supprime
3



simplification. $[ac + (ba + aa)a]^* (aa + ba)$

Exo 2 du TD 8 S et G

5) le langage des mots contenant un nombre pair de a

$$[\varepsilon] = (b^* a b^* a)^* b^*$$

$$[a] = (b^* a b^* a)^* b^* a b^*$$

$$u^{-1}.L = \{w \mid u.w \in L\}$$

$$u \sim v \text{ si } u^{-1}.L = v^{-1}.L$$

$$u \sim v \text{ si } \forall w (u.w \in L \Leftrightarrow v.w \in L) \\ w \in u^{-1}.L \Leftrightarrow w \in v^{-1}.L$$

6) $a^+ b^+$

$$[\varepsilon] = \varepsilon$$

$$\varepsilon^{-1}.L = a^+ b^+$$

$$[a] = a^+$$

$$a^{-1}.L = a^{-1}.(a^+ b^+) = a^{-1}.(a a^+ b^+) = a^+ b^+$$

$$[ab] = a^+ b^+$$

$$(ab)^{-1}.L = b^+$$

$$[b] = \cdot$$

$$b^{-1}.L = \emptyset$$

$$= A^* \setminus \{a^+ + a^+ b^+\} = b^+. (a+b)^* + a^+ b^+ a (a+b)^*$$

$$\begin{aligned}(ab)^{-1}L &= b^{-1} \cdot (a^{-1}L) = b^{-1} \cdot (a^*b^+) \\ &= b^{-1} \cdot (a^*bb^+) = b^x\end{aligned}$$

$$b^{-1}(L_1L_2) = (b^{-1} \cdot L_1)L_2 + b^{-1}L_2 \quad \text{si } L_1 \ni \varepsilon$$

donc

$$\begin{aligned}b^{-1}(a^*bb^+) &= (b^{-1} \cdot a^*)bb^+ + b^{-1} \cdot bb^+ \\ &= \emptyset + b^x\end{aligned}$$

TD 9

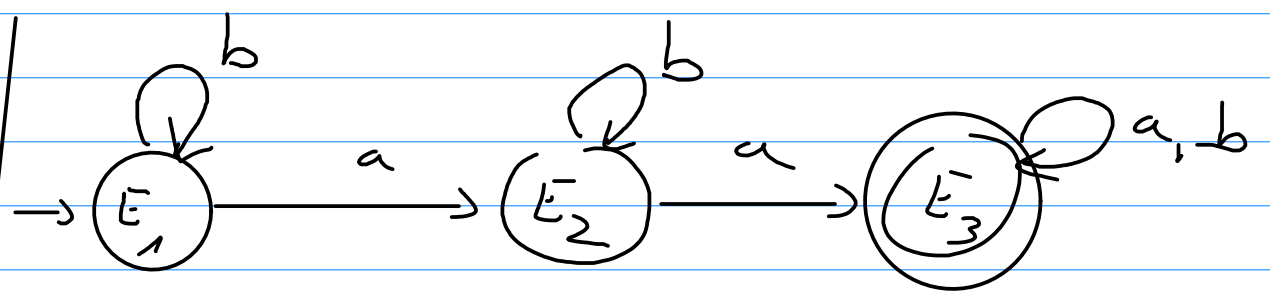
$$L_3 = \{ w \mid |w|_a \geq 2 \} = b^* a b^* a (a+b)^*$$

$$\varepsilon^{-1} L_3 = L_3 = E_1$$

$$a^{-1} L_3 = b^* a (a+b)^* = E_2 \qquad b^{-1} L_3 = L_3 = E_1$$

$$\begin{aligned}
 & a^{-1} L_3 \xrightarrow{a} (aa)^{-1} L_3 = (a+b)^* = E_3 \\
 & a^{-1} L_3 \xrightarrow{b} (ba)^{-1} L_3 = E_2
 \end{aligned}$$

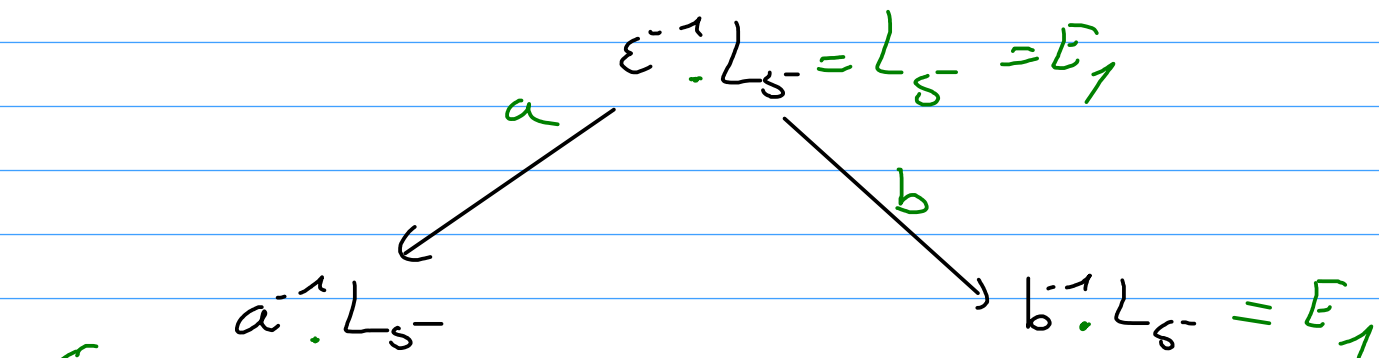
$$\begin{aligned}
 & (aa)^{-1} L_3 = E_3 \\
 & (ba)^{-1} L_3 = E_3
 \end{aligned}$$



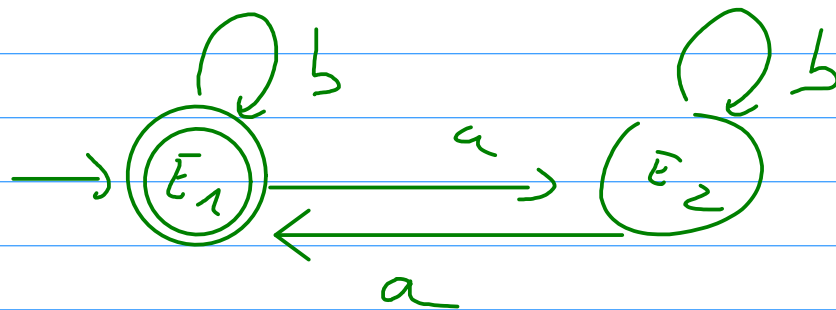
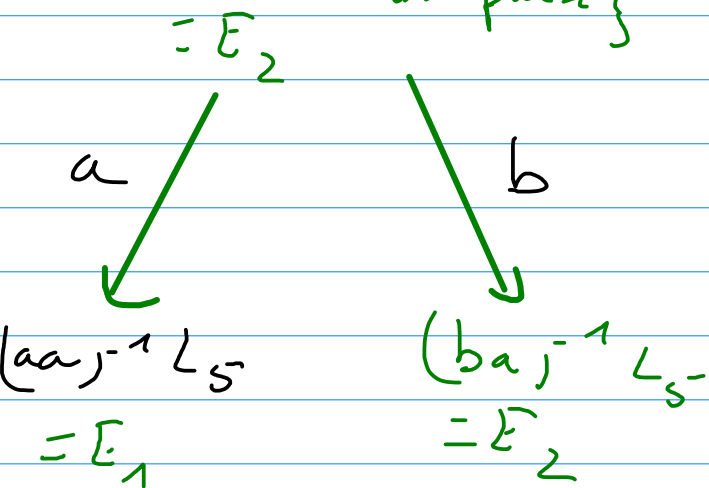
E_i acceptant si $\varepsilon \in E_i$

Automate minimal complet

$$L_S = \{w \in \{a, b\}^* \mid |w|_a \text{ est pair}\}$$



$$\{w \in \{a, b\}^* \mid |w|_a \text{ est impair}\}$$



$\underbrace{aa, \dots, aa}_{\text{nb pair de } a}$
 $\underbrace{ba, \dots, ba}_{\text{nb pair de } a}$

$$(aa)^{-1} \cdot L_S = L_S$$

$$L_4 = \{w \in \{a,b\}^* \mid w \text{ contient le facteur } ab\}$$

$$= (a+b)^* ab (a+b)^*$$

$$\varepsilon^{-1} L_4 = L_4 = E_1$$

